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Identity theory, §8.7

I. The identity predicate is a special predicate, with a special logic

Consider the following logical derivation: Superman can fly. Fs Superman is Clark Kent. ??? So, Clark Kent can fly. Fc

We know the two people are the same, so anything true of one, is true of the other. This is called the Law of the Indiscernibility of Identicals: $(x)(y)[(x=y) \supset (\mathscr{F}x=\mathscr{F}y)]$ '=' is just a special predicate, which we could write as 'Exy' Instead, we introduce a new symbol, '='

The three rules governing Identity (ID)

1)

a=a

For any constant 'a', 'a' is identical to itself. We can add, as a line in any proof, a statement of this form. This only applies to constants.

2)

a=b::b=aIdentity is commutative. This can be applied to variables, as well as to constants.

3)

- Гa
- a=b /Fb

If a=b, then, you may rewrite any formula containing 'a' with 'b' in the place of 'a' throughout.

II. Translation with identity	
Simple ID	
Clark Kent is Superman	c=s
Mary Ann Evans is George Eliot	m=g
Only and except:	
John loves Mary:	Ljm
John only loves Mary:	$Ljm \cdot (x)(Ljx \supset x=m)$
Only john loves Mary:	$Ljm \cdot (x)(Lxm \supset x=j)$
Everyone loves Mary:	$(\mathbf{x})(\mathbf{P}\mathbf{x} \supset \mathbf{L}\mathbf{x}\mathbf{m})$
Everyone except John loves Mary:	$\sim Ljm \cdot (x)[(Px \cdot x \neq j) \supset Lxm]$
Note: ' $x \neq j$ ' is just shorthand for ' $\sim x=j$ '.	
Negation applies to the identity predicate	e, and not to the objects.

ies to the identity predicate, and not to the objects. appi 'B 10

Superlatives:

Degas is a better impressionist than Monet:	Id · Im · Bdm	
Degas is the best impressionist:	$\mathrm{Id} \cdot (\mathbf{x})[(\mathrm{Ix} \cdot \mathbf{x} \neq \mathbf{d}) \supset \mathrm{Bdx}]$	
Bill Gates is the geek with the most money:	$Gg \cdot (x)[(Gx \cdot x \neq g) \supset Mgx]$	
Philadelphia is the nearest major city:	$Mp \cdot (x)[(Mx \cdot x \neq p) \supset Npx]$	
Numerical operators:		
There is only one applicant for the job:	$(\exists x)[Ax \cdot (y)(Ay \supset x=y)]$	
There is at least one applicant for the job:	(∃x)Ax	
There are at least two applicants:	$(\exists x)(\exists y)[(Ax \cdot Ay) \cdot x \neq y]$	
There are exactly two applicants:	$(\exists x)(\exists y)\{[(Ax \cdot Ay) \cdot x \neq y] \cdot (z)[Az \supset (z=x \lor z=y)]\}$	
There is at most one applicant:	$(x)(y)[(Ax \cdot Ay) \supset x=y]$	
There are at most two applicants:	$(x)(y)(z)[(Ax \cdot Ay \cdot Az) \supset (x=y \lor x=z \lor y=z)]$	
Note: this last statement makes no existential commitments!		

III. Definite descriptions

Consider the sentence 'The King of France is bald'. We might translate it as 'Bk'. 'Bk' is false, since there is no King of France. So, '~Bk' should be true, since it's the negation of a false statement. That means that 'It's false that the King of France is bald' is true. Which seems to imply that the King of France has hair. In fact, we want both 'The King of France is bald' and 'The King of France is not bald' to be false. So, we had better translate the sentence differently.

'The King of France' is a definite description.

It refers to one specific object without using a name.

There are two ways to refer: by name or by description (E.g. the person who, the thing that)

Bertrand Russell's analysis for definite descriptions:

Both sentences 'The King of France is bald' and 'The King of France is not bald' are false, due to a false presupposition. They're really complex statements, involving a definite description.

'The King of France is bald' entails three simpler expressions:

A. There is a King of France.

- B. There is only one King of France.
- C. That thing is bald.

So, the proposition is false because clause A is false. 'The King of France is not bald' is also false, for the same reason. The negation only affects the third clause. The first is still the same, and still false.

Another example: The country called a sub-continent is India.

1. There is a country called a sub-continent.

2. There is only one such country.

3. That country is identical with India.

So, we translate as: $(\exists x) \{ (Cx \cdot Sx) \cdot (y) [(Cy \cdot Sy) \supset y=x] \cdot x=I \}$

Another example: 'The author of Waverly was a genius' becomes ' $(\exists x) \{Wx \cdot (y)[Wy \supset y=x] \cdot Gx\}$ '

IV. More examples

1. Everything is identical with itself. (x)x=x2. Nothing is distinct from itself. $(x) \sim x = x$ 3. The discoverer of Polonium is Polish. (Dx, Px) $(\exists x) \{ Dx \cdot (y) (Dy \supset y=x) \cdot Px \}$ 4. At most two persons invented the airplane. (Px, Ix) $(\mathbf{x})(\mathbf{y})(\mathbf{z})[(\mathbf{P}\mathbf{x} \cdot \mathbf{I}\mathbf{x} \cdot \mathbf{P}\mathbf{y} \cdot \mathbf{I}\mathbf{y} \cdot \mathbf{P}\mathbf{z} \cdot \mathbf{I}\mathbf{z}) \supset (\mathbf{x}=\mathbf{y} \lor \mathbf{x}=\mathbf{z} \lor \mathbf{y}=\mathbf{z})]$ 5. There is exactly one dollar bill in my wallet. (Dx, Wx) $(\exists x) \{ (Dx \cdot Wx) \cdot (y) [(Dy \cdot Wy) \supset y=x] \}$ 6. Adriana is a better dancer than Rene. (a, r, Dx, Bxy: x is better than y) $Da \cdot Dr \cdot Bar$ 7. Adriana is the best dancer. $Da \cdot (x)[(Dx \cdot \ x=a) \supset Bax]$ 8. There are exactly two dancers better than Adriana. $(\exists x)(\exists y)\{(Dx \cdot Bxa \cdot Dy \cdot Bya \cdot \neg x=y) \cdot (z)[(Dz \cdot Bza) \supset (z=x \lor z=y)]\}$

V. Exercises A. Translate the following sentences into predicate logic using '='.

- 1. Everything is identical with something.
- 2. There are prime numbers. (Px, Nx)
- 3. Two is the only even prime number. (t, Ex, Px, Nx)
- 4. There is exactly one even prime number.
- 5. There are at least two odd prime numbers. (Ox, Px, Nx)
- 6. All prime numbers are odd except the number two.
- 7. The murderer was Colonel Mustard. (m, Mx)
- 8. There is at least one dancer better than Rene.
- 9. Rene is the worst dancer.
- 10. Goliath is the tallest human. (g, Hx, Txy: x is taller than y)

VI. Derivations in identity theory

Consider the original problem: Superman can fly.

Superman is Clark Kent.

 \therefore Clark Kent can fly.

1. Fs	
2. s=c	/ Fc
3. Fc	1, 2, ID
QED	

Using the commutative Id rule:

1. a=b ⊃ j=k	
2. b=a	
3. Fj	/ Fk
4. a=b	2, Id
5. j=k	1, 4, MP
6. Fk	3, 5, Id
QED	

To derive the negation of an identity statement, you must use (IP):

4, 5, Conj

1. Rm 2. ~ Rj / m≠j 3. m=j 4. Rj 5. Rj · ~ Rj 6. m≠j QED Using the reflexive Id rule: 1. $(\mathbf{x})(\sim \mathbf{G}\mathbf{x} \supset \mathbf{x} \neq \mathbf{d})$ / Gd 2. ~Gd AIP 3. $\sim Gd \supset d \neq d$ 1, UI 4. d=d Id 5. d≠d 3, 2, MP

7. Gd QED

Proof with existential conclusion:

6. $d=d \cdot d \neq d$

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1. Rab
2. (\exists x) \sim Rxb
                                  /(\exists x) \sim x = a
3. \sim Rcb
                                             2, EI
             4. c=a
                                             AIP
             5. Rcb
                                                         1, Id
             6. Rcb \cdot \sim Rcb
                                             5, 3, Conj
7. \sim c=a
                                             4-6, IP
8. (\exists x) \sim x = a
                                             7, EG
QED
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A final example: The Faulkner scholar at Swarthmore is very learned. Therefore, all Faulkner scholars at Swarthmore are very learned.

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1. (\exists x) {(Sx \cdot Fx) \cdot (y)[(Sy \cdot Fy) \supset x=y]} · Lx
                                                                 /(\mathbf{x})[(\mathbf{Sx} \cdot \mathbf{Fx}) \supset \mathbf{Lx}]
            2. \sim(x)[(Sx · Fx) \supset Lx]
                                                                             AIP
            3. (\exists x) \sim [(Sx \cdot Fx) \supset Lx]
                                                                             2, CQ
            4. ~ [(Sa \cdot Fa) \supset La]
                                                                             3, EI
            5. ~ [~ (Sa \cdot Fa) \vee La]
                                                                             4, Impl
                                                                             5, DM, DN
            6. (Sa \cdot Fa) \cdot ~La
            7. {(Sb · Fb) · (y)[(Sy · Fy) \supset b=y]} · Lb
                                                                             1, EI (to b)
            8. (y)[(Sy \cdot Fy) \supset b=y]
                                                                                         7, Simp, Com, Simp
            9. (Sa \cdot Fa) \supset b=a
                                                                             8, UI (to a)
            10. Sa · Fa
                                                                             6, Simp
            11. b=a
                                                                             9, 10, MP
            12. Lb
                                                                             7, Simp
            13. La
                                                                                         12, 11, Id
            14. ~La
                                                                             6, Com, Simp
            15. La · ~ La
                                                                             13, 14, Conj
16. (x)[(Sx \cdot Fx) \supset Lx]
                                                                             2-15, IP
QED
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VII. Exercises B. Derive the conclusions of each of the following arguments.

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1)
1. (x)(Dx \supset Ex)
2. Da
3. a=b
                                       / Eb
2)
1. (x)(Ax \supset Bx)
2. ~Bf
3. Ae
                                       / f≠e
3)
1. (x)(Hx \supset Jx)
2. (x)(Kx \supset Lx)
3. Hd \cdot Kc
4. c=d
                                       / Jc · Ld
4)
1. (x)(y)(x=y)
2. (x)Mxx
                                       / Mab
5)
1. (x)[(\exists y)Kxy \supset (\exists z)Kzx]
2. (\exists x)(Kxg \cdot x=b)
                                       /(∃z)Kzb
6)
1. (∃x)Hx
2. (x)(y)[(Hx \cdot Hy) \supset x=y]
                                                /(\exists x)[Hx \cdot (y)(Hy \supset x=y)]
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Solutions may vary.

VIII. Solutions

Answers to Exercises A. 1. $(x)(\exists y)x=y$ 2. $(\exists x)(Px \cdot Nx)$ 3. $Et \cdot Pt \cdot Nt \cdot (x)[(Ex \cdot Px \cdot Nx) \supset x=t]$ 4. $(\exists x)\{(Ex \cdot Px \cdot Nx) \cdot (y)[(Ey \cdot Py \cdot Ny) \supset y=x]\}$ 5. $(\exists x)(\exists y)(Ox \cdot Px \cdot Nx \cdot Oy \cdot Py \cdot Ny \cdot \neg x=y)$ 6. $(x)[(Px \cdot Nx \cdot \neg x=t) \supset Ox]$ 7. $(\exists x)[Mx \cdot (y)(My \supset y=x) \cdot x=m]$ 8. $Dr \cdot (\exists x)(Dx \cdot Bxr)$ 9. $Dr \cdot (x)[(Dx \cdot \neg x=r) \supset Bxr]$ 10. $Hg \cdot (x)[(Hx \cdot \neg x=g) \supset Tgx]$ Or, using definite descriptions: $(\exists x)\{(y)[(Hy \cdot \neg y=x) \supset Txy] \cdot (z)\{(y)[(Hy \cdot \neg y=z) \supset Tzy] \supset z=x\} \cdot x=g\}$ A solution to Exercise B.6: 1. (∃x)Hx 2. $(x)(y)[(Hx \cdot Hy) \supset x=y]$ $/(\exists x)[Hx \cdot (y)(Hy \supset x=y)]$ 3. Ha 1, EI 4. $\sim (\exists x)[Hx \cdot (y)(Hy \supset x=y)]$ AIP 5. $(x) \sim [Hx \cdot (y)(Hy \supset x=y)]$ 4, CQ 5, DM 6. $(x)[\sim Hx \lor \sim (y)(Hy \supset x=y)]$ 7. ~Ha \lor ~(y)(Hy \supset a=y) 6, UI 7, 3, DN, DS 8. \sim (y)(Hy \supset a=y) 9. $(\exists y) \sim (Hy \supset a=y)$ 8, CQ 10. ~(Hb \supset a=b) 9, EI 11. ~(~Hb ∨ a=b) 10, Impl 12. Hb · ~a=b 11, DM, DN 13.Hb 12, Simp 12, Com Simp 14.~a=b 15. (y)[(Ha · Hy) \supset a=y] 2, UI 16. (Ha · Hb) \supset a=b 15, UI 17. Ha · Hb 3, 13, Conj 18. a=b 16, 17, MP 19. $a=b \cdot a=b$ 18, 14, Conj 4-19, IP 20. $(\exists x)[Hx \cdot (y)(Hy \supset x=y)]$