

Philosophy 240: Symbolic Logic

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Fall 2015

Class #40 - Functions

Business

- Functions today
 - Translation and derivation
- Second-order logic on Wednesday
 - Just translation
 - We won't do inferences
- Friday: review, maybe reflect
- Final is Monday the 14th, 7pm
 - Review session on Sunday the 13th, 4pm
 - Here
- Compensatory Exams
 - Up to two
 - Respond to email which I will send
- Course Evaluations
- What else?



A Motivating Argument for Functions

1. No odd numbers are even.
 2. One is odd.
 3. One is the square of one.
- So, not all square numbers are even.

- We can regiment into **F**.

1. $(\forall x)(Ox \supset \sim Ex)$
2. Oo
3. $(\exists x)[Sxo \cdot (\forall y)(Syo \supset y=x) \cdot x=o]$
 $/ \sim(\forall x)[(\exists y)Sxy \supset Ex]$

- But, there is a more efficient and fecund option.
- Take ‘the square of x’ as a function.

Functions

- A small extension of **F** introduces functors to represent functions.
- A function is a special kind of relation.
- An object may bear the same relation to various different objects
 - ▶ Laa, Lab, Lac...
 - ▶ Gab, Gcb, Gdb...
 - ▶ I am taller than lots of things and younger than lots of things.
 - ▶ I love several things.
- A function associates a given object (or given objects) with exactly one object.
 - ▶ An n -place relation in which one place of the relation is unique for given n -tuples of the other places.
 - ▶ It takes one or more arguments and returns a single output, called its range.
 - one-place functions take one argument
 - two-place functions take two arguments
 - n -place functions take n arguments

Functions Are All Over

- Mathematics
 - linear functions
 - exponential functions
 - periodic functions
 - quadratic functions
 - trigonometric functions.
- Science
 - force is a function of mass and acceleration
 - momentum is a function of mass and velocity
 - Your genetic code is a function of the codes of your biological parents.
- Logic
 - semantics for **PL**: truth functions
- Natural language
 - the (biological) father of
 - the (biological) teacher of

Functions and Uniqueness

- Functions have a unique, determinate value for any given input.
- Any human being has one biological father
 - Putting genetic engineering aside
- Most relations are not functions
 - A person can **love** or **be loved** (**know** or **be known**) by many people
 - Many different things can be **between** x and z.
- These are not functions
 - the parents of a
 - the classes that a and b share
 - the square root of x

Some Functions and Their Logical Representations

- the father of: $f(x)$
- the successor of: $g(x)$
- the sum of: $f(x,y)$
- the teacher of: $g(x_1 \dots x_n)$
 - Given no team teaching!
- The truth value of the conjunction of a and b: $f(a,b)$

Vocabulary of FF

- Capital letters A...Z, used as predicates
- Lower case letters
 - ▶ a, b, c, d, e, i, j, k...u are used as constants.
 - ▶ **f, g, and h are used as functors.**
 - ▶ v, w, x, y, z are used as variables.
- Five connectives: \sim , \bullet , \vee , \supset , \equiv
- Quantifiers: \exists , \forall
- Punctuation: $()$, $[]$, $\{\}$

N-Tuples

- A functor term is a functor symbol followed by an n-tuple of singular terms.
- An **n-tuple of singular terms** is an ordered series of terms.
 - ▶ Singular terms: constants, variables, or functor terms
 - ▶ 'single', 'double', 'triple', 'quadruple', etc. are n-tuples.
 - ▶ Functions can take any number of arguments.
- We use n-tuples in the semantics of relational predicates.
 - ▶ In the metalanguage
 - ▶ Often: $\langle a, b, c \rangle$
- We will represent n-tuples in our object language merely by listing the terms separated by commas.
- Some n-tuples
 - ▶ a, b
 - ▶ $a, a, f(a)$
 - ▶ $x, y, b, d, f(x), f(a, b, f(x))$
 - ▶ a

Functor Terms

- If α is an n-tuple of singular terms, then the following are all **functor terms**:
 - ▶ $f(\alpha)$
 - ▶ $g(\alpha)$
 - ▶ $h(\alpha)$
- Note that an n-tuple of terms can include functor terms.
 - ▶ 'Functor term' is defined recursively.
 - ▶ We allow composition of functions.
- We can refer to the grandfather of x using just the functions for father, e.g. ' $f(x)$ ', and mother, e.g. ' $g(x)$ ':
 - ▶ $f(f(x))$
 - ▶ $f(g(x))$
- Composition of mathematical functions
 - ▶ Take ' $h(x)$ ' to represent the square of x
 - ▶ ' $h(h(h(x)))$ ' represents the eighth power of x , i.e. $((x^2)^2)^2$.

Formation Rules for Wffs of FF

1. An n-place predicate followed by n singular terms (constants, variables, **or functor terms**) is a wff.
2. For any variable β , if α is a wff that does not contain either ' $(\exists\beta)$ ' or ' $(\forall\beta)$ ', then ' $(\exists\beta)\alpha$ ' and ' $(\forall\beta)\alpha$ ' are wffs.
3. If α is a wff, so is $\sim\alpha$.
4. If α and β are wffs, then so are:
 - ▶ $(\alpha \cdot \beta)$
 - ▶ $(\alpha \vee \beta)$
 - ▶ $(\alpha \supset \beta)$
 - ▶ $(\alpha \equiv \beta)$
5. These are the only ways to make wffs.

The scope and binding rules are the same for **FF** as they were for **M** and **F**.

FF: Semantics

- The semantics for **FF** are basically the same as for **F**.
- We insert an interpretation of function symbols.
 - ▶ Step 1. Specify a set to serve as a domain of interpretation, or domain of quantification.
 - ▶ Step 2. Assign a member of the domain to each constant.
 - ▶ **Step 3. Assign a (metalinguistic) function with arguments and ranges in the domain to each function symbol.**
 - ▶ Step 4. Assign some set of objects in the domain to each one-place predicate; assign sets of ordered n-tuples to each relational predicate.
 - ▶ Step 5. Use the customary truth tables for the interpretation of the connectives.
- The function assigned in Step 3 will be a function in the meta-language used to interpret the function in the object language.
 - ▶ Remember, functions are just a kind of relation.
 - ▶ They don't need any new bells or whistles.

Translations Into FF

- Translation key:
 - Lxy : x loves y
 - $f(x)$: the father of x
 - $g(x)$: the mother of x
- Olaf loves his mother.
 - $\text{Log}(o)$
- Olaf loves his grandmothers.
 - $\text{Log}(g(o)) \bullet \text{Log}(f(o))$
- No one is his/her own mother.
 - $(\forall x)\sim x=g(x)$
- No one is her/his own grandmother.
 - $(\forall x)[\sim x=g(f(x)) \bullet \sim x=f(f(x))]$

Functions and Mathematics

- Many simple concepts in arithmetic are functions.
 - addition
 - multiplication
 - least common multiple
- The most fundamental function in mathematics is the successor function.
- All other mathematical functions can be defined in terms of successor and other basic concepts.
- All of arithmetic can be developed from five basic axioms.
- **Peano's Axioms for Arithmetic**
 - P1: 0 is a number
 - P2: The successor (x') of every number (x) is a number
 - P3: 0 is not the successor of any number
 - P4: If $x'=y'$ then $x=y$
 - P5: If P is a property that may (or may not) hold for any number, and if
 - i. 0 has P ; and
 - ii. for any x , if x has P then x' has P ;then all numbers have P .

Peano's Axioms, Regimented

Key: a: zero; Nx: x is a number; f(x): the successor of x

P1: 0 is a number

P1. Na

P2: The successor (x') of every number (x) is a number

P2. $(\forall x)(Nx \supset Nf(x))$

P3: 0 is not the successor of any number

P3. $\sim(\exists x)(Nx \cdot f(x)=a)$

P4: If $x'=y'$ then $x=y$

P4. $(\forall x)(\forall y)[(Nx \cdot Ny) \supset (f(x)=f(y) \supset x=y)]$

P5: If P is a property that may hold for any number, and if

- i. 0 has P; and
- ii. for any x, if x has P then x' has P; then all numbers have P.

P5. $\{Pa \cdot (\forall x)[(Nx \cdot Px) \supset Pf(x)]\} \supset (\forall x)(Nx \supset Px)$

Some Number-Theoretic Statements

▶ Key:

- ▶ o : one
- ▶ $f(x)$: the successor of x
- ▶ $f(x, y)$: the product of x and y
- ▶ Ex : x is even
- ▶ Ox : x is odd
- ▶ Px : x is prime

1. One is not the successor of any number.

- ▶ $(\forall x)(Nx \supset \sim f(x)=o)$

2. If the product of a pair of numbers is odd, then the product of the successors of those numbers is even.

- ▶ $(\forall x)(\forall y)\{(Nx \cdot Ny) \supset [Of(x, y) \supset Ef(f(x), f(y))]\}$

3. There are no prime numbers such that their product is prime.

- ▶ $\sim(\exists x)(\exists y)[Nx \cdot Px \cdot Ny \cdot Py \cdot Pf(x, y)]$

Derivations Using Functions

- No new rules
- Functions are singular terms.
- A functor can be either a constant or a variable.
 - ▶ It depends on what the arguments of the function are.
- We can UI to a variable, or a function of a variable, or any complex function all of whose arguments are variables.
- For UG, if the arguments of a function are all variables, then we are free to use UG over the variables in that function.
 - ▶ If the arguments contain any constants, then we can not use UG.
 - ▶ The restrictions on UG continue to hold for variables which are arguments of a function.
 - CP and IP
 - If a constant is present when the variable is introduced
- For EI, we must continue always to instantiate to a new term.
 - ▶ A functor is not a new term if any of its arguments, or any of the arguments of any of its sub-functors, have already appeared in the derivation.

The Motivating Argument

1. No odd numbers are even.
 2. One is odd.
 3. One is the square of one.
- So, not all square numbers are even.

1. $(\forall x)(Ox \supset \sim Ex)$
 2. Oo
 3. $o=f(o)$
- $\therefore \sim(\forall x)Ef(x)$

More Derivations

1. $(\forall x)[Ax \supset Bxf(x)]$

2. $(\exists x)Af(x) \quad / \quad (\exists x)Bf(x)f(f(x))$

1. $\sim(\exists x)Cx \quad / \quad (\forall x)\sim Cf(x, g(x))$

1. $(\forall x)\{(Nx \cdot Gxt) \supset (\exists y)(\exists z)[Py \cdot Pz \cdot x=f(y, z)]\}$

2. $Nb \cdot Gbt \quad / \quad (\exists x)(\exists y)(\exists z)[Nx \cdot Py \cdot Pz \cdot x=f(y, z)]$