

Philosophy 240

Symbolic Logic

Russell Marcus
Hamilton College
Fall 2013

Class #20: On Axiomatic Systems

Hunter's System PS

With axiom schemata

- Three axiom schemata:

$$\text{PS1: } \alpha \supset (\beta \supset \alpha)$$

$$\text{PS2: } (\alpha \supset (\beta \supset \gamma)) \supset ((\alpha \supset \beta) \supset (\alpha \supset \gamma))$$

$$\text{PS3: } (\sim \alpha \supset \sim \beta) \supset (\beta \supset \alpha)$$

Any formula of the language PS of any of the forms PS1, PS2, or PS3 is an axiom of PS.

– Infinitely many axioms of PS

- One rule of inference, modus ponens.

If α and β are formulas, β is a consequence in system PS of α and $\alpha \supset \beta$.

$$\alpha, (\alpha \supset \beta) \vdash_{\text{PS}} \beta$$

Hunter's System PS

With a substitution rule

- Three axioms

$$\text{PS1*}: P \supset (Q \supset P)$$

$$\text{PS2*}: (P \supset (Q \supset R)) \supset ((P \supset Q) \supset (P \supset R))$$

$$\text{PS3*}: (\sim P \supset \sim Q) \supset (Q \supset P)$$

- Two rules of inference

$$\text{MP: } \alpha, (\alpha \supset \beta) \vdash_{\text{PS}} \beta$$

Substitution: For any wff α , and any wff β appearing as a sub-formula in wff δ : If $\vdash_{\text{PS}} \delta$, then $\vdash_{\text{PS}} \delta_{\alpha}^{\beta}$
– where ' δ_{α}^{β} ' is the result of substituting ' α ' for ' β ' throughout.

Hunter's PS and Our PL

- Hunter's axiomatic system is provably equivalent to our more-familiar natural deduction system PL.
 - ▶ A metalogical result (the deduction theorem) allows us to move between the two kinds of systems.
- Frege's original *Begriffsschrift* used a Hilbert-style axiomatic system, with a different axiomatization.
 - ▶ See the appendices in Richard Mendelsohn's *The Philosophy of Gottlob Frege*, for a translation of the *Begriffsschrift* into modern notation.
- Natural deduction systems like PS in *What Follows* are due largely to Gerhard Gentzen's work in the 1930s and 1940s.
 - ▶ There is a paper on the history of natural deduction by Pelletier which might be worth a look for a term paper.
- Completeness:
 - ▶ Frege's axiomatization, Hunter's PS, and standard systems of natural deduction are all complete.
 - All logical truths are provable within the system.
- One's choice of system is thus arbitrary among the various complete systems.

A Derivation in PS

PS1: $\alpha \supset (\beta \supset \alpha)$

PS2: $(\alpha \supset (\beta \supset \gamma)) \supset ((\alpha \supset \beta) \supset (\alpha \supset \gamma))$

PS3: $(\sim\alpha \supset \sim\beta) \supset (\beta \supset \alpha)$

■ A1: $\vdash_{PS} P \supset P$

1. $P \supset ((P \supset P) \supset P)$

PS1

2. $(P \supset ((P \supset P) \supset P)) \supset ((P \supset (P \supset P)) \supset (P \supset P))$

PS2

3. $(P \supset (P \supset P)) \supset (P \supset P)$

MP, 1, 2

4. $P \supset (P \supset P)$

PS1

5. $P \supset P$

3, 4, MP

■ QED

A Further Derivation in PS

PS1: $\alpha \supset (\beta \supset \alpha)$

PS2: $(\alpha \supset (\beta \supset \gamma)) \supset ((\alpha \supset \beta) \supset (\alpha \supset \gamma))$

PS3: $(\sim\alpha \supset \sim\beta) \supset (\beta \supset \alpha)$

A1: $P \supset P$

■ A2: $\vdash_{PS} \sim P \supset (P \supset P)$

1. $P \supset P$

Ex A1

2. $(P \supset P) \supset (\sim P \supset (P \supset P))$

PS 1

3. $\sim P \supset (P \supset P)$

MP, 1, 2

■ QED

A Longer Proof in PS

■ QED