

# **Philosophy 240: Symbolic Logic**

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Class 40 - Functions

# Final Exam

Thursday

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7pm - 10pm

# A Motivating Argument for Functions

1. No odd numbers are even.
  2. One is odd.
  3. One is the square of one.
- So, not all square numbers are even.

- We can regiment into **F**.
  1.  $(\forall x)(Ox \supset \sim Ex)$
  2.  $Oo$
  3.  $(\exists x)[Sxo \cdot (\forall y)(Syo \supset y=x) \cdot x=o]$   
 $/ \sim(\forall x)[(Sx \cdot Nx) \supset Ex]$
- But, there is a more efficient, and more fecund, option.
- Take ‘the square of x’ as a function.

# Functions

- A small extension of **F** introduces functors to represent functions.
- A function takes one or more arguments and returns a single output, its range.
- Mathematics
  - linear functions
  - exponential functions
  - periodic functions
  - quadratic functions
  - trigonometric functions.
- Science
  - force is a function of mass and acceleration
  - momentum is a function of mass and velocity
  - genetic code
- Logic
  - semantics for **PL**
- Natural language
  - the father of
  - the teacher of
- One-place functions take one argument, two-place functions take two arguments, n-place functions take n arguments.

# Some Functions and Their Logical Representations

- the father of:  $f(x)$
- the successor of:  $g(x)$
- the sum of:  $f(x,y)$
- the teacher of:  $g(x_1 \dots x_n)$
- These are not functions:
  - the parents of  $a$
  - the classes that  $a$  and  $b$  share
  - the square root of  $x$

# Vocabulary of FF

- Capital letters A...Z, used as predicates
- Lower case letters
  - ▶ a, b, c, d, e, i, j, k...u are used as constants.
  - ▶ **f, g, and h are used as functors.**
  - ▶ v, w, x, y, z are used as variables.
- Five connectives:  $\sim$ ,  $\bullet$ ,  $\vee$ ,  $\supset$ ,  $\equiv$
- Quantifiers:  $\exists$ ,  $\forall$
- Punctuation:  $()$ ,  $[]$ ,  $\{\}$

# N-Tuples

- A functor term is a functor symbol followed by an n-tuple of singular terms.
- An **n-tuple of singular terms** is an ordered series of terms.
  - ▶ Singular terms: constants, variables, or functor terms
  - ▶ 'single', 'double', 'triple', 'quadruple', etc. are n-tuples.
  - ▶ Functions can take any number of arguments.
- We use n-tuples in the semantics of relational predicates.
- Often:  $\langle a, b, c \rangle$
- We will represent n-tuples merely by listing the terms separated by commas.
- Some n-tuples
  - ▶  $a, b$
  - ▶  $a, a, f(a)$
  - ▶  $x, y, b, d, f(x), f(a, b, f(x))$
  - ▶  $a$

# Functor Terms

- If  $\alpha$  is an n-tuple of singular terms, then the following are all **functor terms**:
  - $f(\alpha)$
  - $g(\alpha)$
  - $h(\alpha)$
- Note that an n-tuple of terms can include functor terms.
- 'Functor term' is defined recursively, which allows for composition of functions.
- For example, one can refer to the grandfather of  $x$ , using just the functions for father, e.g. ' $f(x)$ ', and mother, e.g. ' $g(x)$ ':
  - $f(f(x))$
  - $f(g(x))$
- Composition of mathematical functions
  - Take ' $h(x)$ ' to represent the square of  $x$
  - ' $h(h(h(x)))$ ' represents the eighth power of  $x$ , i.e.  $((x^2)^2)^2$ .



# Formation Rules for Wffs of FF

1. An  $n$ -place predicate followed by  $n$  singular terms (constants, variables, **or functor terms**) is a wff.
2. For any variable  $\beta$ , if  $\alpha$  is a wff that does not contain either ' $(\exists\beta)$ ' or ' $(\forall\beta)$ ', then ' $(\exists\beta)\alpha$ ' and ' $(\forall\beta)\alpha$ ' are wffs.
3. If  $\alpha$  is a wff, so is  $\sim\alpha$ .
4. If  $\alpha$  and  $\beta$  are wffs, then so are:
  - ▶  $(\alpha \cdot \beta)$
  - ▶  $(\alpha \vee \beta)$
  - ▶  $(\alpha \supset \beta)$
  - ▶  $(\alpha \equiv \beta)$
5. These are the only ways to make wffs.

The scope and binding rules are the same for **FF** as they were for **M** and **F**.

# FF: Semantics

- The semantics for **FF** are basically the same as for **F**.
- We insert an interpretation of function symbols.
  - ▶ Step 1. Specify a set to serve as a domain of interpretation, or domain of quantification.
  - ▶ Step 2. Assign a member of the domain to each constant.
  - ▶ **Step 3. Assign a function with arguments and ranges in the domain to each function symbol.**
  - ▶ Step 4. Assign some set of objects in the domain to each one-place predicate; assign sets of ordered n-tuples to each relational predicate.
  - ▶ Step 5. Use the customary truth tables for the interpretation of the connectives.
- The function assigned in Step 3 will be a function in the meta-language used to interpret the function in the object language.

# Translations Into FF

- Translation key:
  - $Lxy$ : x loves y
  - $f(x)$ :the father of x
  - $g(x)$ :the mother of x
- Olaf loves his mother
  - $Log(o)$
- Olaf loves his grandmothers
  - $Log(g(o)) \cdot Log(f(o))$
- No one is his/her own mother
  - $(\forall x) \sim x=g(x)$

# Functions and Mathematics

- Many simple concepts in arithmetic are functions: addition, multiplication, least common multiple.
- The most fundamental function in mathematics is the successor function.
- All other mathematical functions can be defined in terms of successor and other basic concepts.
- All of arithmetic can be developed from five basic axioms, called the Peano axioms.
- Peano's Axioms for Arithmetic
  - P1: 0 is a number
  - P2: The successor ( $x'$ ) of every number ( $x$ ) is a number
  - P3: 0 is not the successor of any number
  - P4: If  $x'=y'$  then  $x=y$
  - P5: If  $P$  is a property that may (or may not) hold for any number, and if
    - i. 0 has  $P$ ; and
    - ii. for any  $x$ , if  $x$  has  $P$  then  $x'$  has  $P$ ;then all numbers have  $P$ .

# Peano's Axioms, Regimented

Key: a: zero; Nx: x is a number; f(x): the successor of x

P1. Na

P2.  $(\forall x)(Nx \supset Nf(x))$

P3.  $\sim(\exists x)(Nx \cdot f(x)=a)$

P4.  $(\forall x)(\forall y)[(Nx \cdot Ny) \supset (f(x)=f(y) \supset x=y)]$

P5.  $\{Pa \cdot (\forall x)[(Nx \cdot Px) \supset Pf(x)]\} \supset (\forall x)(Nx \supset Px)$

# Some Number-Theoretic Statements

▶ Key:

- ▶  $o$ : one
- ▶  $f(x)$ : the successor of  $x$
- ▶  $f(x, y)$ : the product of  $x$  and  $y$
- ▶  $Ex$ :  $x$  is even
- ▶  $Ox$ :  $x$  is odd
- ▶  $Px$ :  $x$  is prime

1. One is not the successor of any number.

- ▶  $(\forall x)(Nx \supset \sim f(x)=o)$

2. If the product of a pair of numbers is odd, then the product of the successors of those numbers is even.

- ▶  $(\forall x)(\forall y)\{(Nx \cdot Ny) \supset [Of(x, y) \supset Ef(f(x), f(y))]\}$

3. There are no prime numbers such that their product is prime.

- ▶  $\sim(\exists x)(\exists y)[Nx \cdot Px \cdot Ny \cdot Py \cdot Pf(x, y)]$

# Derivations Using Functions

- No new rules
- Functions act like simple terms.
- A functor can be either a constant or a variable.
  - It depends on what the arguments of the function are.
- We can UI to a variable, or a function of a variable, or any complex function all of whose arguments are variables.
- For UG, if the arguments of a function are all variables, then we are free to use UG over the variables in that function.
  - If the arguments contain any constants, then we can not use UG.
- For EI, we must continue always to instantiate to a new term.
  - A functor is not a new term if any of its arguments, or any of the arguments of any of its sub-functors, have already appeared in the derivation.

# The Motivating Argument

1. No odd numbers are even.
  2. One is odd.
  3. One is the square of one.
- So, not all square numbers are even.

$$1. (\forall x)(Ox \supset \sim Ex)$$

$$2. Oo$$

$$3. o=f(o)$$

$$/ \sim(\forall x)Ef(x)$$



# More Derivations

1.  $(\forall x)[Ax \supset Bxf(x)]$

2.  $(\exists x)Af(x) \quad / \quad (\exists x)Bf(x)f(f(x))$

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1.  $\sim(\exists x)Cx \quad / \quad (\forall x)\sim Cf(x, g(x))$

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1.  $(\forall x)\{(Nx \cdot Gxt) \supset (\exists y)(\exists z)[Py \cdot Pz \cdot x=f(y, z)]\}$

2.  $Nb \cdot Gbt \quad / \quad (\exists x)(\exists y)(\exists z)[Nx \cdot Py \cdot Pz \cdot x=f(y, z)]$